

## Dynamic Wavelength and Bandwidth Allocation for Power Saving in TWDM PON

Man Soo Han

Dept. of Information and Communications Eng., Mokpo National Univ.,  
Jeonnam, Republic of Korea  
[mshan@mokpo.ac.kr](mailto:mshan@mokpo.ac.kr)

**Abstract.** This paper proposes a dynamic bandwidth and bandwidth allocation (DWBA) algorithm for time and wavelength division multiplexing passive optical networks (TWDM PONs). The proposed algorithm saves power by decreasing the number of active wavelengths of a TWDM PON system.

**Keywords:** TWDM, PON, DWBA, power saving.

### 1 Introduction

Time and wavelength division multiplexing passive optical network (TWDM PON) is a solution for next generation PON stage 2 (NG-PON2) that supports 40 Gbps in the upstream and 10 up to 40 Gbps in the downstream [1]. A TWDM PON consists of an optical line termination (OLT) and a number of optical network units (ONUs). Also multiple wavelengths are used in both upstream and downstream directions. In the downstream direction, broadcasting is used to transmit packets from an OLT to ONUs. In the upstream direction, the OLT performs a dynamic wavelength and bandwidth allocation (DWBA) to dynamically assign an upstream wavelength and to dynamically allocate a non-overlapping transmission slot to each ONU.

Access networks consume about 80% of the energy consumed in the Internet [2]. A TWDM PON system has potential for saving energy since it uses multiple wavelengths and many ONUs. Most of DWBA algorithms of TWDM PON systems are focused on quality-of-service (QoS) such as a mean delay and a pack loss rate [3][4][5]. The energy consumption is proportional to the number of active wavelengths in a TWDM PON system. To save energy, DWBA algorithms must reduce the number of active wavelengths. Energy saving DWBA algorithm was introduced in [2]. However, the TWDM PON system of [2] is based on EPON (Ethernet PON). The TWDM PON technology [1][6] is different from the EPON technology in many points.

In this paper, we propose two DWBA algorithms for a TWDM PON system which is compliant to the TWDM PON recommendations. The first algorithm is a DWBA for QoS (DAQ) whose object is to obtain high QoS by fully using the wavelengths. The second algorithm is a DWBA for power saving (DAP) whose aim is to decrease the active wavelengths to save power.

## 2 Power Saving DWBA

A TWDM PON system consists of a single OLT and N ONUs. The OLT has four transmitters and four receivers. Four downstream wavelengths and four upstream wavelengths are used in the system. Each ONU has a single tunable transmitter and a single tunable receiver. An ONU dynamically uses one of the upstream wavelengths and one of the downstream wavelengths.

To support QoS, three service classes are used in the system. In the TWDM PON technology, a service class is known as a transmission container (T-CONT) type. Each ONU has three queues for T-CONT types 2, 3 and 4, respectively. Since a static bandwidth allocation is used for the T-CONT type 1, we do not consider the T-CONT type 1 in this paper. A packet arrived from a user side to an ONU is saved one of the queues based on its T-CONT type. Each queue has a unique Allocation Identifier (AllocID).

Every operation of TWDM PON is synchronized with a frame duration (FD) which is fixed at 125  $\mu$ s. Each queue reports its queue length as a request to the OLT. The OLT saves the requests to perform a DWBA for an upcoming FD. The OLT performs a DWBA in every FD to assign an upstream wavelength and grant a transmission slot to each queue. The total sum of the transmission slots is constrained by the duration of an FD.

In this paper, we assume that an ONU does not change its upstream wavelength during an FD. When a transmitter changes its wavelength, it requires a switching time in which no packet can be transmitted. The switching time decreases the bandwidth utilization. Another side effect is that a transmission overhead increases. When an ONU starts a new transmission, it needs a guard time which is included in the transmission overhead. If an ONU changes its upstream wavelength, a new guard time has to be inserted. For simplicity, we assume that an ONU uses only one upstream wavelength in an FD. For the purpose of comparison, we first develop a new DWBA algorithm for QoS (DAQ) whose sole aim is to support high QoS such as a low mean delay. Then we introduce a new DWBA algorithm for power saving (DAP) whose object is to decrease power consumption.

In [5], a DWBA algorithm is developed for a TWDM PON system with a single T-CONT type. The main object of the DWBA algorithm of [5] is QoS. To get high QoS, an ONU selects an upstream wavelength which has the largest remaining bandwidth in [5]. Extending the DWBA algorithm of [5], we develop a DAQ algorithm for the TWDM PON system with multiple T-CONT types.

Let queue(j) be a queue having AllocID j. Suppose request(j) and grant(j) are the request and the grant of queue(j), respectively. The queue(j) has two parameters S(j) and A(j): S(j) is the service interval in unit of FD and A(j) is the maximum service bytes that can be allocated during S(j). The queue(j) has two counters T(j) and V(j). The counter T(j) is decreased by 1 in each FD and then recharged to S(j) when it has expired. The counter V(j) denotes the remaining service bytes during S(j) and is decreased by the grant amount grant(j). The counter V(j) is reset to A(j) when the counter T(j) has expired.

Let W(i) be the upstream wavelength of ONU i where  $W(i) = 1, \dots, 4$ . If an upstream wavelength of ONU i is 4, we write as  $W(i) = 4$ . If no upstream wavelength is assigned to ONU i, then we write as  $W(i) = 0$ . Assume F(k) is the remaining

bandwidth of an upstream wavelength  $k$ . At the beginning of each FD,  $F(k)$  is charged to its maximum value. The bandwidth  $F(k)$  is decreased by the grant amount during a DWBA process.

The first stage of the DAQ is the allocation of the upstream wavelength. Since ONU  $i$  uses a single upstream wavelength during an FD, all queues of ONU  $i$  have to use the same upstream wavelength during an FD. Suppose  $queue(j)$  is a queue of ONU  $i$  and  $grant(j)$  has not been allocated. If the upstream wavelength of ONU  $i$  has been allocated, i.e.,  $W(i) > 0$ ,  $queue(j)$  has to use the upstream wavelength  $W(i)$ . If the upstream wavelength of ONU  $i$  has not been assigned, i.e.,  $W(i) = 0$ , the OLT assigns a tentative wavelength,  $C$ , to ONU  $i$ . The tentative wavelength  $C$  is given by the index of the maximum of  $F(k)$ ,  $k=1, \dots, 4$ . For instance,  $F(2)$  is the largest one then we have  $C = 2$ . The tentative wavelength  $C$  will be finally allocated to ONU  $i$  only when  $grant(j) > 0$ . That is  $W(i) = C$  only if  $grant(j) > 0$ . Otherwise,  $W(i) = 0$ .

The second stage of the DAQ is the allocation of bandwidth. The  $grant(j)$  is calculated by

$$grant(j) = \min(request(j), V(j), F(k)) \quad (1)$$

where  $k = W(i)$  if  $W(i) > 0$ , otherwise  $k = C$ . If  $grant(j) > 0$  and  $W(i) = 0$ , then the tentative wavelength  $C$  is allocated to ONU  $i$ , i.e.,  $W(i) = C$ . If  $grant(j) = 0$  and  $W(i) = 0$ , we have  $W(i) = 0$ . At the end of the calculation, each of  $request(j)$ ,  $V(j)$  and  $F(k)$  is decreased by the amount of  $grant(j)$ .

For the service fairness among ONUs, the DAQ algorithm changes the starting ONU number in a round robin manner. For example, if the DAQ started from ONU 4 at the previous FD, the DAQ starts from ONU 5 at the current FD. The DAQ algorithm uses a static priority among T-CONT types. First, the DAQ algorithm is performed for the queues of T-CONT type 2. Then the algorithm is executed for the queues of T-CONT type 3 with the remaining resources. Finally the algorithm is processed for the queues of T-CONT type 4 with the remaining resources.

Since the DAQ prefers the wavelength having the largest remaining bandwidth, the upstream wavelengths are fully used in general. Since the power consumption is proportional to the number of active upstream wavelengths, the DAQ is not efficient for energy saving. We propose a DAP algorithm whose main purpose is power saving in this subsection. To save energy, the DAP decreases the number of active upstream wavelengths.

The DAP uses the same parameters and counters of the DAQ. The proposed DAP algorithm comprises of three stages. The first stage of the DAP is the calculation of an estimated number of upstream wavelengths,  $E$ . Suppose  $queue(j)$  is a queue of ONU  $i$  and  $grant(j)$  has not been allocated. Let  $G(j)$  be the minimum of  $request(j)$  and  $V(j)$ . That is  $G(j) = \min(request(j), V(j))$ . Then  $grant(j)$  must be less than or equal to  $G(j)$ . Then the estimated number  $E$  is given by

$$E = \lceil \sum_j G(j) / M \rceil \quad (2)$$

where the variable  $M$  is the maximum of  $F(k)$  and the maximum of  $E$  is 4.

The second stage of the DAP is the allocation of the upstream wavelength which is similar to that of the DAQ. The number of upstream wavelengths is limited by  $E$ . If the upstream wavelength of ONU  $i$  has been allocated, i.e.,  $W(i) > 0$ ,  $queue(j)$  uses the upstream wavelength  $W(i)$ . If the upstream wavelength of ONU  $i$  has not been assigned, i.e.,  $W(i) = 0$ , the OLT assigns a tentative wavelength,  $C$ , to ONU  $i$ . The tentative wavelength  $C$  is given by the index of the maximum of  $F(k)$ ,  $k=1, \dots, E$ .

The third stage is the bandwidth allocation. The  $\text{grant}(j)$  is calculated by

$$\text{grant}(j) = \min(G(j), F(k)) \quad (3)$$

where  $k = W(i)$  if  $W(i) > 0$ , otherwise  $k = C$ . If  $\text{grant}(j) > 0$  and  $W(i) = 0$ , then the tentative wavelength  $C$  is allocated to ONU  $i$ , i.e.,  $W(i) = C$ . If  $\text{grant}(j) = 0$  and  $W(i) = 0$ , we have  $W(i) = 0$ . At the end of the calculation, each of  $\text{request}(j)$ ,  $V(j)$  and  $F(k)$  is decreased by the amount of  $\text{grant}(j)$ .

## References

1. ITU-T Rec. G.989.2 Draft Version, 40-Gigabit-capable passive optical networks (NG-PON2): Physical media dependent (PMD) layer specification, (2013)
2. Dixit A., Lannoo B., Colle D., Pickavert M., Demeester P.: Novel DBA algorithm for Energy Efficiency in TWDM-PONs, 39th European Conference and Exhibition on Optical Communication (ECOC), pp. 1--3, (2013)
3. Dixit A., Lannoo B., Colle D., Pickavert M., Demeester P.: Dynamic Bandwidth Allocation with Optimal Wavelength Switching in TWDM-PONs, 15th International Conference on Transparent Optical Networks (ICTON), pp. 1--4, (2013)
4. Zhang J., Ansari N.: Scheduling Hybrid WDM/TDM Passive Optical Networks With Nonzero Laser Tuning Time, IEEE/ACM Transactions on Networking, vol. 19, no. 4, pp. 1014--1027, (2011)
5. Han M. S.: Performance Evaluation of Dynamic Bandwidth Allocation Algorithm for TWDM PON, Journal of Convergence Information Technology, vol. 8, no. 13, pp.282--288, (2013)
6. ITU-T Rec. G.987.3 Rev.2, 10-Gigabit-capable passive optical networks (XG-PON): Transmission convergence (TC) specifications, (2010)
7. Han M. S., Yoo H., Lee D. S.: Development of Efficient Dynamic Bandwidth Allocation algorithm for XGPON, ETRI Journal, vol. 3, no. 1, pp. 18--26, (2013)